



Analysis and Design of Algorithms

Lecture 2



Analysis of Algorithms I

Dr. Mohamed Loey

Lecturer, Faculty of Computers and Information
Benha University
Egypt

Table of Contents

Analysis of Algorithms

Time complexity

Asymptotic Notations

Big O Notation

Growth Orders

Problems

Analysis of Algorithms

- **Analysis of Algorithms** is the determination of the amount of **time, storage** and/or other **resources** necessary to execute them.
- Analyzing algorithms is called **Asymptotic Analysis**
- **Asymptotic Analysis** evaluate the performance of an algorithm

Time complexity

Time complexity

- ❑ **time complexity** of an algorithm quantifies the amount of **time** taken by an algorithm

- ❑ We can have three cases to analyze an algorithm:
 - 1) Worst Case
 - 2) Average Case
 - 3) Best Case

Time complexity

□ Assume the below algorithm using Python code:

```
def search(arr, x):  
    for i in range(len(arr)):  
        if arr[i] == x:  
            return i+1  
    return -1
```

Time complexity

- **Worst Case Analysis:** In the worst case analysis, we calculate upper bound on running time of an algorithm.

```
def search(arr, x):  
    for i in range(len(arr)):  
        if arr[i] == x:  
            return i+1  
    return -1
```

Time complexity

- ❑ **Worst Case Analysis:** the case that causes maximum number of operations to be executed.
- ❑ For Linear Search, the worst case happens when the element to be searched is not present in the array. (example : search for number 8)

2	3	5	4	1	7	6
----------	----------	----------	----------	----------	----------	----------

Time complexity

- ❑ **Worst Case Analysis:** When x is not present, the `search()` functions compares it with all the elements of `arr` one by one.

```
def search(arr, x):  
    for i in range(len(arr)):  
        if arr[i] == x:  
            return i+1  
    return -1
```

Time complexity

- The worst case time complexity of linear search would be $O(n)$.

```
def search(arr, x):  
    for i in range(len(arr)):  
        if arr[i] == x:  
            return i+1  
    return -1
```

Time complexity

- **Average Case Analysis:** we take all possible inputs and calculate computing time for all of the inputs.

```
def search(arr, x):  
    for i in range(len(arr)):  
        if arr[i] == x:  
            return i+1  
    return -1
```

Time complexity

- **Best Case Analysis:** calculate lower bound on running time of an algorithm.

```
def search(arr, x):  
    for i in range(len(arr)):  
        if arr[i] == x:  
            return i+1  
    return -1
```

Time complexity

- The best case time complexity of linear search would be $O(1)$.

```
def search(arr, x):  
    for i in range(len(arr)):  
        if arr[i] == x:  
            return i+1  
    return -1
```

Time complexity

- ❑ **Best Case Analysis:** the case that causes minimum number of operations to be executed.
- ❑ For Linear Search, the best case occurs when x is present at the first location. (example : search for number 2)
- ❑ So time complexity in the best case would be $\Theta(1)$

2	3	5	4	1	7	6
----------	----------	----------	----------	----------	----------	----------

Time complexity

- ❑ Most of the times, we do **worst case** analysis to analyze algorithms.
- ❑ The **average case** analysis is not easy to do in most of the practical cases and it is rarely done.
- ❑ The **Best case** analysis is bogus. Guaranteeing a lower bound on an algorithm doesn't provide any information.

Asymptotic Notations

- 1) **Big O Notation**: is an Asymptotic Notation for the worst case.
- 2) **Ω Notation** (omega notation): is an Asymptotic Notation for the best case.
- 3) **Θ Notation** (theta notation) : is an Asymptotic Notation for the worst case and the best case.

Big O Notation

Big O Notation

1) $O(1)$

- Time complexity of a function (or set of statements) is considered as $O(1)$ if it doesn't contain loop, recursion and call to any other non-constant time function. For example `swap()` function has $O(1)$ time complexity.

```
def swap(s1, s2):  
    return s2, s1
```

Big O Notation

- A loop or recursion that runs a constant number of times is also considered as $O(1)$. For example the following loop is $O(1)$.

```
# c is constant  
c=4  
for i in range(1,c):  
    #some O(1) expressions  
    #print(i)
```

2) $O(n)$

- Time Complexity of a loop is considered as $O(n)$ if the loop variables is incremented / decremented by a constant amount. For example the following loop statements have $O(n)$ time complexity.

2) $O(n)$

```
# n is variable  
# c is increment  
for i in range(1, n, c):  
    #some  $O(1)$  expressions  
    print(i)
```

2) $O(n)$

□ Another Example:

```
// Here c is a positive integer constant
for (int i = 1; i <= n; i += c) {
    // some O(1) expressions
}

for (int i = n; i > 0; i -= c) {
    // some O(1) expressions
}
```

3) $O(n^c)$

- Time complexity of nested loops is equal to the number of times the innermost statement is executed. For example the following loop statements have $O(n^2)$ time complexity

3) $O(n^2)$

```
# n is variable  
# c is increment  
for i in range(1,n,c):  
    #some O(1) expressions  
    for j in range(1,n,c):  
        #some O(1) expressions  
        print(i,j)
```


□ Another Example

```
for (int i = 1; i <=n; i += c) {  
    for (int j = 1; j <=n; j += c) {  
        // some O(1) expressions  
    }  
}  
  
for (int i = n; i > 0; i -= c) {  
    for (int j = i+1; j <=n; j += c) {  
        // some O(1) expressions  
    }  
}
```

Big O Notation

4) $O(\text{Log}n)$

- Time Complexity of a loop is considered as $O(\text{Log}n)$ if the loop variables is divided / multiplied by a constant amount.

```
# n is variable  
# c is constant  
i=2  
while i<=n:  
    print(i)  
    i=i*c
```

4) $O(\text{Log}n)$

➤ Another Example

```
for (int i = 1; i <=n; i *= c) {  
    // some  $O(1)$  expressions  
}  
for (int i = n; i > 0; i /= c) {  
    // some  $O(1)$  expressions  
}
```

5) $O(\text{LogLog}n)$

- Time Complexity of a loop is considered as $O(\text{LogLog}n)$ if the loop variables is reduced / increased exponentially by a constant.

```
# n is variable  
# c is constant  
i=2  
while i<=n:  
    print(i)  
    i=i**c
```

5) $O(\text{LogLog}n)$

➤ Another Example

```
// Here c is a constant greater than 1
for (int i = 2; i <=n; i = pow(i, c)) {
    // some O(1) expressions
}
//Here fun is sqrt or cuberoot or any other constant root
for (int i = n; i > 0; i = fun(i)) {
    // some O(1) expressions
}
```

- How to combine time complexities of consecutive loops?

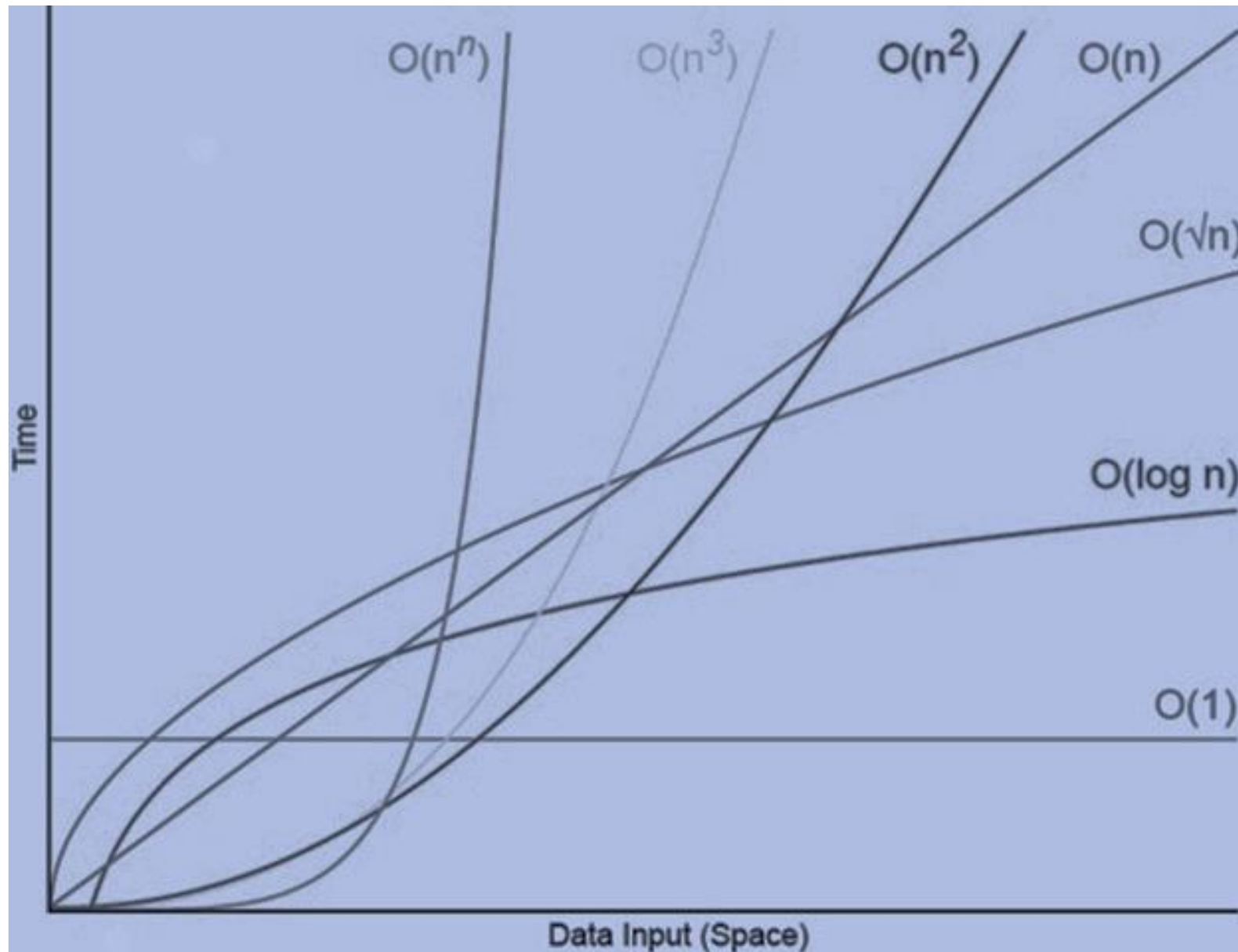
```
# n, k is variable  
  
for i in range(n):  
    print(i)  
for j in range(m):  
    print(j)
```

- Time complexity of above code is $O(n) + O(m)$ which is $O(n+m)$

Growth Orders

n	O(1)	O(log(n))	O(n)	O(nlog(n))	O(N²)	O(2ⁿ)	O(n!)
1	1	0	1	1	1	2	1
8	1	3	8	24	64	256	40x10³
30	1	5	30	150	900	10x10⁹	210x10³²
500	1	9	500	4500	25x10⁴	3x10¹⁵⁰	1x10¹¹³⁴
1000	1	10	1000	10x10³	1x10⁶	1x10³⁰¹	4x10²⁵⁶⁷
16x10³	1	14	16x10³	224x10³	256x10⁶	-	-
1x10⁵	1	17	1x10⁵	17x10⁵	10x10⁹	-	-

Growth Orders



Growth Orders

Length of Input (N)	Worst Accepted Algorithm
≤ 10	$O(N!), O(N^6)$
≤ 15	$O(2^N * N^2)$
≤ 20	$O(2^N * N)$
≤ 100	$O(N^4)$
≤ 400	$O(N^3)$
$\leq 2K$	$O(N^2 * \log N)$
$\leq 10K$	$O(N^2)$
$\leq 1M$	$O(N * \log N)$
$\leq 100M$	$O(N), O(\log N), O(1)$

Problems

- Find the complexity of the below program:

```
function(int n)
{
    if (n==1)
        return;
    for (int i=1; i<=n; i++)
    {
        for (int j=1; j<=n; j++)
        {
            printf("*");
            break;
        }
    }
}
```

Problems

- **Solution:** Time Complexity $O(n)$.
Even though the inner loop is bounded by n , but due to break statement it is executing only once.

```
function(int n)
{
    if (n==1)
        return;
    for (int i=1; i<=n; i++)
    {
        // Inner loop executes only one
        // time due to break statement.
        for (int j=1; j<=n; j++)
        {
            printf("*");
            break;
        }
    }
}
```

Problems

□ Find the complexity of the below program:

```
void function(int n)
{
    int count = 0;
    for (int i=n/2; i<=n; i++)
        for (int j=1; j+n/2<=n; j = j++)
            for (int k=1; k<=n; k = k * 2)
                count++;
}
```

Problems

□ Solution:

Time

$O(n^2 \log n)$

```
void function(int n)
{
    int count = 0;

    // outer loop executes n/2 times
    for (int i=n/2; i<=n; i++)

        // middle loop executes n/2 times
        for (int j=1; j+n/2<=n; j = j++)

            // inner loop executes logn times
            for (int k=1; k<=n; k = k * 2)
                count++;
}
```

Problems

□ Find the complexity of the below program:

```
void function(int n)
{
    int count = 0;
    for (int i=n/2; i<=n; i++)
        for (int j=1; j<=n; j = 2 * j)
            for (int k=1; k<=n; k = k * 2)
                count++;
}
```

Problems

□ Solution:

Time

$O(n \log^2 n)$

```
void function(int n)
{
    int count = 0;
    for (int i=n/2; i<=n; i++)

        // Executes  $O(\log n)$  times
        for (int j=1; j<=n; j = 2 * j)

            // Executes  $O(\log n)$  times
            for (int k=1; k<=n; k = k * 2)
                count++;
}
```

Problems

- Find the complexity of the below program:

```
void function(int n)
{
    int count = 0;
    for (int i=0; i<n; i++)
        for (int j=i; j< i*i; j++)
            if (j%i == 0)
                for (int k=0; k<j; k++)
                    printf("*");
}
```


Problems

□ Solution:

Time $O(n^5)$

```
void function(int n)
{
    int count = 0;

    // executes n times
    for (int i=0; i<n; i++)

        // executes O(n*n) times.
        for (int j=i; j< i*i; j++)
            if (j%i == 0)
            {
                // executes j times = O(n*n) times
                for (int k=0; k<j; k++)
                    printf("*");
            }
}
```

Contact Me



**THANKS FOR
YOUR TIME**

